

# CS 2461

## Lab- Week 2

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### Today....

- Quick review of data representation and operations on bits
- Review Transistor circuits (gates)

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## Binary Representation Summary

- Every storage locations stores a finite sequence of bits
  - 8-bit, 16-bit, 32-bit etc.
- The same bit string can mean different things depending on how the program wants to look at it....based on data representation used

Address	7	6	5	4	3	2	1	0	
35	1	0	0	0	0	0	0	1	Unsigned: +129
36	1	0	0	0	0	1	1	1	2C: -127
37	1	1	1	0	0	0	0	1	2C: 109
38	0	1	1	0	1	1	0	1	ASCII: 'm'

If string at address 38 is defined as type `int` in C then value = 109  
 If defined as `char` then value='m'

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## Arithmetic and Logic Operations

- Arithmetic:
  - Addition
  - Subtract (negative number and add to second number)
  - Shift – left shift by one position is multiplying by 2; right shift is division by 2
    - Shift left twice = multiply by  $2^2 = 4$
  - Multiplication....
- Logic operators
  - AND, OR, NOT,..... Define using truth tables
  - Bitwise operations – apply logical operator at each bit position

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## Shifting Bit Fields

	7	6	5	4	3	2	1	0
Original Pattern x	0	1	1	0	1	0	1	1
X << 1 – Left Shift by 1	1	1	0	1	0	1	1	0
X << 2 – Left Shift by 2	1	0	1	0	1	1	0	0
Original Pattern x	1	1	1	0	1	0	1	1
X >> 1 – Shift Right (logical) by 1	0	1	1	1	0	1	0	1
X >>> 1 – Shift Right (arithmetic) by 1	1	1	1	1	0	1	0	1

- Shift Left:
  - Move all #'s to the left, fill in empty spots with a 0
- Shift Right (2 kinds):
  - shift right logical (SRL) >>
    - shift 0's in from the left
  - shift right arithmetic (SRA) >>>
    - replicate the sign bit, (**very useful for sign extension!**)

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## Dose of reality: Finite Width and Overflow

- Integers have infinite width
  - There are an infinite number of them
- **Hardware integers have finite (architecture defined) width**
  - Limited by hardware circuits themselves
  - 64-bit these days ( $2^{64}$  integers):
  - LC3 integers are 16-bit ( $2^{16}$  or ~64,000)
- **Overflow**: when operation result is outside type's range
  - Example: 15 + 1 with 4-bit integers (16 needs 5 bits, 10000)

$$\begin{array}{r}
 \text{overflow} \\
 \text{(carry-out)} \\
 + \quad \text{1111}_{(5)} \\
 + \quad \text{0001}_{(4)} \\
 \hline
 \text{10000}_{(16)}
 \end{array}$$

*Problem: using 4-bit representation the sum is 0!!*

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## Arithmetic Overflow - Summary

- For **unsigned** numbers
  - Any addition that produces an 'extra bit' is a problem
- For **2C signed** numbers
  - Sometimes addition or subtraction produce an extra bit – **this is not necessarily a problem.**
  - *Overflow if Signs of both operands are the same AND the sign of the sum is different*
  - Arithmetic overflow can occur when you are adding 2 positive or 2 negative numbers – in this case if the sign of the result is different from the sign of the addends you have an arithmetic overflow
    - (this is the key to determining overflow condition in 2C)
  - CPU architectures today, use 2C representation

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## Sign Extension

- Suppose we have a number which is stored in a four bit register and we want to add this number to a number stored in a eight bit register
- We have a device (an 8-bit adder) which will do the addition and it is designed to add two 8 bit numbers
  - SW Analogy: Calling a function with the correct (type, number) arguments
- Therefore extend 4-bit number to 8-bit.... How ?
- Suppose we just pad 0's to the left:
  - 4 bit 0100 (decimal 4) becomes 0000 0100 which is decimal 4
  - 4 bit 1100 (decimal -4) becomes 0000 1100 which is decimal 12...wrong!
- Solution: replicate Most Significant bit (pad MSB to the left)
  - 4 bit 0100 (decimal 4) becomes 0000 0100 which is still decimal 4
  - 4 bit 1100 (decimal -4) becomes 1111 1100 which is still decimal -4

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## Bitwise Logical Operations

- View n-bit field as a collection of n logical values

- *Apply operation to each bit independently*

- Bitwise AND: useful for clearing bits

- AND with zero = 0
- AND with one = no change

AND  $\begin{array}{r} 11000101 \\ 00001111 \\ \hline \end{array}$

- Bitwise OR: useful for setting bits

- OR with zero = no change
- OR with one = 1

$\begin{array}{r} 00000101 \\ 11000101 \\ \hline \end{array}$   
OR

- Computers don't support individual bits as a data type

- Just use least significant bit of n-bit integer
- Integers are generally more useful

$\begin{array}{r} 00001111 \\ \hline 11001111 \end{array}$

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## Bitwise Operations – Group Exercises.....

- Let A,B,C,D,F be *any* 8 bit 2's complement numbers
    - i.e., think of A as type `int` in a C program
- What is  $C = A \text{ AND } 00000001$ 
    - How many different values can C have? What are they and when do they occur?
    - What property of A is determined by this "statement"?
  - What is  $D = (A \gg 7) \text{ AND } 00000001$  (*right shift operator*)
    - How many different values can D have? When do they occur?
    - What property of A is determined by this "statement"?
  - What is  $F = (A \text{ OR } B) \text{ AND } 00000001$ 
    - How many different value can F have? When does each value occur?
    - What property (of A,B) is determined by this "statement"?

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## Review....Transistor Circuits and Logic gates

- Transistor acts as a voltage controlled switch
  - Send 0 or 1 to transistor Gate: switch closes or opens
- Two types of Transistors used in our circuits:
  - p-type = switch closes if gate input =0, open if input =1
  - n-type = switch closes if gate input =1, open if input =0
- Circuit Output = voltage measured at some location in the circuit
  - 1 if there is a positive voltage, and 0 if no voltage

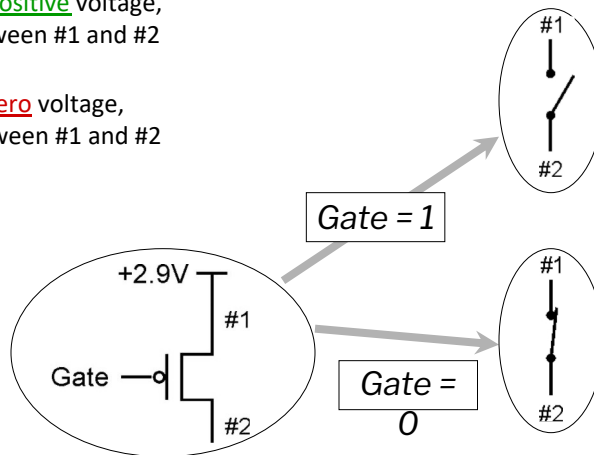
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## Abstraction: Simplified view of p-type MOS Transistor

### ▪ p-type

- when Gate has positive voltage, open circuit between #1 and #2 (switch open)
- when Gate has zero voltage, short circuit between #1 and #2 (switch closed)



**Important: For p-type, Terminal #1 must be connected to Voltage Source.**

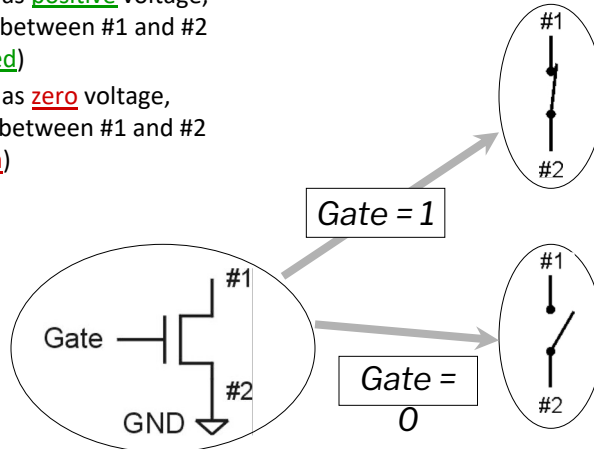
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## Abstraction: Simplified view of n-type MOS Transistor

▪ n-type *complementary* to p-type

- when Gate has positive voltage, short circuit between #1 and #2 (switch closed)
- when Gate has zero voltage, open circuit between #1 and #2 (switch open)

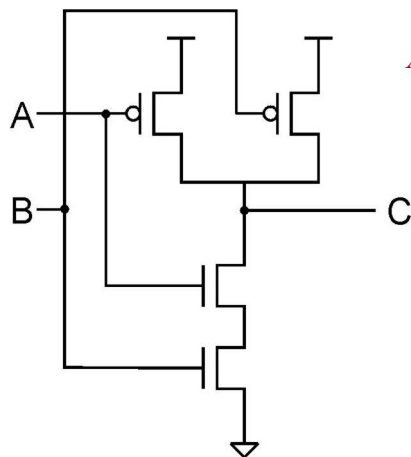


**Important: For n-type, Terminal #2 must be connected to Ground (0V).**

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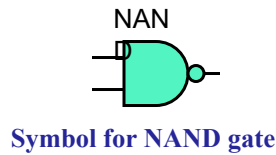
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## NAND Gate – NOT (AND): $C = NOT (A AND B)$



*So how to build an AND gate?  
 $A AND B = NOT (NOT (A AND B))$*

*Use NAND and NOT  
 Send output of NAND to  
 Input of NOT gate*



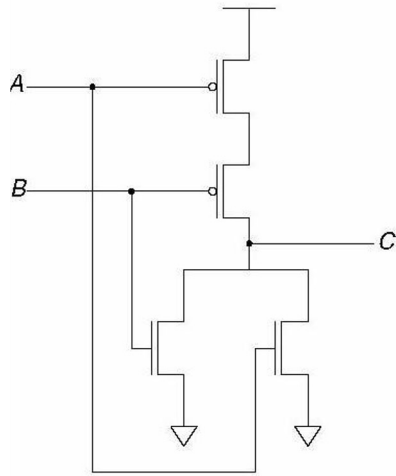
Truth Table

A	B	C
0	0	1
0	1	1
1	0	1
1	1	0

Note: Parallel structure on top, serial on bottom.

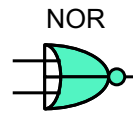
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### Solution to Question – Group Exercise



- 1. Derive truth table for this circuit, Inputs= A,B Output =C
- 2. What function is implemented ?

NOR gate = NOT (A OR B)

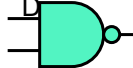


### Basic Logic Gates - Symbols

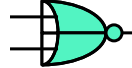
NOT/INV



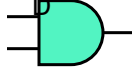
NAN



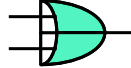
NOR



AN




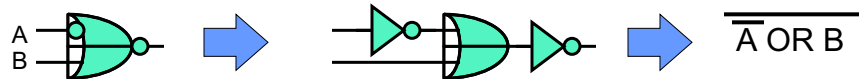
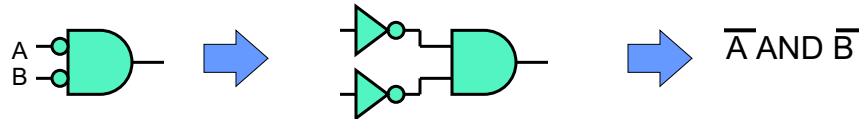
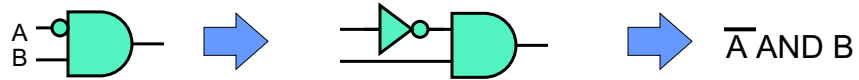
OR





### Shorthand for Inverting Signals

- Invert a signal by adding either
  - a  before/after a gate
  - a "bar" over letter

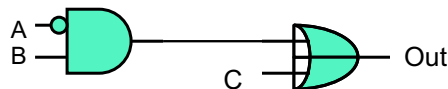


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### Example: Your first combinational circuit

- Combinational logic circuits ~ propositional logic statements
- Use gates to implement the logic operators ( 'functions' )
  - *No necessity to show the circuit using transistors since each gate corresponds to an implementation using transistors*
- Output = ( (NOT A) AND B ) OR C
- Need one AND gate and one OR gate (and one NOT gate/invertor)

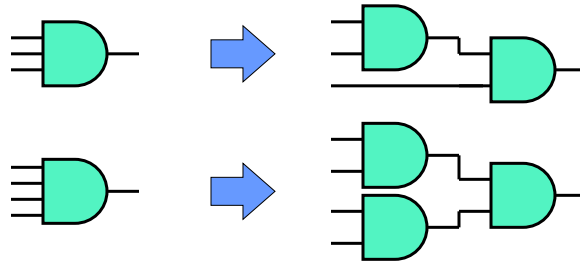


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### More than 2 Inputs? Arbitrary Functions?

- AND/OR can take any number of inputs
  - AND = 1 if all inputs are 1
  - OR = 1 if any input is 1 (0 if all inputs are 0)
- Implementation
  - Multiple two-input gates or single CMOS circuit

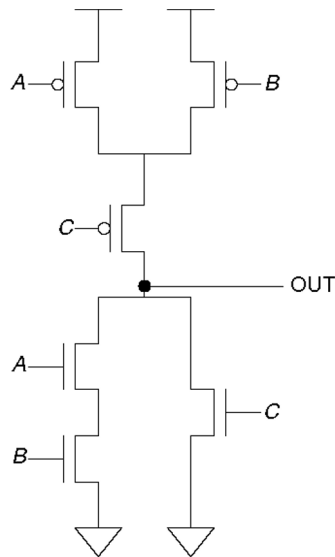


- Can implement arbitrary boolean functions as a gate
  - More complex n- and p- networks

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### Exercise....Truth table for circuit



A	B	C	OUT
0	0	0	
0	0	1	
0	1	0	
0	1	1	
1	0	0	
1	0	1	
1	1	0	
1	1	1	

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